

**IN THE UNITED STATES  
PATENT AND TRADEMARK OFFICE**

**PATENT APPLICATION**

5 Applicant(s): Kedar Sharadchandra Namjoshi  
Case: 2  
Serial No.: 10/614,618  
Filing Date: July 7, 2003  
10 Group: 2193  
Examiner: Tuan A. Vu

Title: Method and Apparatus for Reducing a Program Size While Maintaining  
Branching Time Properties and Automated Checking of Such Reduced Programs  
15

---

APPEAL BRIEF

20 Mail Stop Appeal Brief - Patents  
Commissioner for Patents  
P.O. Box 1450  
Alexandria, VA 22313-1450

25 Sir:

Applicant hereby appeals the non-final rejection dated October 31, 2007, of  
claims 1 through 26 of the above-identified patent application. Applicant hereby submits a  
Request to Reinstate Appeal. An Appeal Brief was submitted on August 14, 2007. Applicant  
30 respectfully requests that the Examiner allow this application to proceed to the Appeal Board

REAL PARTY IN INTEREST

The present application is assigned to Lucent Technologies Inc., as evidenced by  
an assignment recorded on July 7, 2003 in the United States Patent and Trademark Office at Reel  
35 014318, Frame 0212. The assignee, Lucent Technologies Inc., is the real party in interest.

RELATED APPEALS AND INTERFERENCES

There are no related appeals or interferences.

STATUS OF CLAIMS

5            Claims 1 through 26 are presently pending in the above-identified patent application. The specification is objected to for incorporating subject matter by reference. Claims 1-26 are rejected under 35 U.S.C. §112, second paragraph, as being indefinite for failing to particularly point out and distinctly claim the subject matter which applicant regards as the invention. Claims 1-4, 12-17 and 25-26 are rejected under 35 U.S.C. §102(b) as being  
10 anticipated by BenAri (Ben-Ari et al., "The Temporal Logic of Branching Time," Acta Information 20, 6, 207-26(1983)); claims 5-7 and 18-20 are rejected under 35 U.S.C. §103(a) as being unpatentable over BenAri. Claims 8-11 were indicated to contain allowable subject matter. Claims 21-24 were not explicitly addressed by the Examiner but are believed to similarly contain allowable subject matter, since they are substantially similar to claims 8-11.  
15 Claims 1 and 14 are being appealed.

STATUS OF AMENDMENTS

There have been no amendments filed subsequent to the final rejection.

20            SUMMARY OF CLAIMED SUBJECT MATTER

Independent claim 1 is directed to a method (FIG. 1: 200) for reducing a program, M, that preserves at least one branching time property, f (page 3, lines 20-29), comprising the steps of: forming a product of said program, M and said branching time property, f, expressed as an automaton, A (page 6, lines 18-22; FIG. 2, step 210); obtaining an abstract domain containing  
25 a set of abstract values to generalize possible states of said program and abstract relations that relate said program states to said abstract domain (page 6, line 18, to page 7, line 3; FIG. 2, step 220); and computing an abstract program with a reduced number of states and an altered version

of said branching time property, f, using said product (page 3, lines 20-29; FIG. 2, step 240).

Independent claim 14 is directed to a system for reducing a program, M, that preserves at least one branching time property, f (page 3, lines 20-29), comprising: a memory; and a processor operatively coupled to said memory, said processor configured to: form a  
5 product of said program, M and said branching time property, f, expressed as an automaton, A (page 6, lines 18-22; FIG. 2, step 210); obtain an abstract domain containing a set of abstract values to generalize possible states of said program and abstract relations that relate said program states to said abstract domain (page 6, line 18, to page 7, line 3; FIG. 2, step 220); and compute  
10 an abstract program with a reduced number of states and an altered version of said branching time property, f, using said product (page 3, lines 20-29; FIG. 2, step 240).

#### STATEMENT OF GROUNDS OF REJECTION TO BE REVIEWED ON APPEAL

The specification is objected to for incorporating subject matter by reference. Claims 1-26 are rejected under 35 U.S.C. §112, second paragraph, as being indefinite for failing  
15 to particularly point out and distinctly claim the subject matter which applicant regards as the invention. Claims 1-4, 12-17 and 25-26 are rejected under 35 U.S.C. §102(b) as being anticipated by BenAri (Ben-Ari et al., "The Temporal Logic of Branching Time," Acta Information 20, 6, 207-26(1983)); claims 5-7 and 18-20 are rejected under 35 U.S.C. §103(a) as being unpatentable over BenAri. Claims 8-11 were indicated to contain allowable subject  
20 matter. Claims 21-24 were not explicitly addressed by the Examiner but are believed to similarly contain allowable subject matter, since they are substantially similar to claims 8-11.

#### ARGUMENT

##### Specification Objection

25 The Examiner has requested hard copies of any NPL documents that may help the Examiner to better interpret and understand what is novel. Appellant again note that only non-essential material provided for background information has been incorporated by reference.

Appellant would be more than happy to provide a hard copy of any document that would help the Examiner better understand the invention. It is difficult, however, for Appellant to ascertain which portions of the invention the Examiner is struggling with. As set forth in the above Summary section, for example, Appellant maintains that every claim limitation is fully supported  
5 and described within the specification.

Section 112 Rejection

Claims 1-26 were rejected under 35 U.S.C. §112, second paragraph, as being indefinite for failing to particularly point out and distinctly claim the subject matter which applicant regards as the invention. The Examiner asserts that the claims omit essential steps.  
10 The *only* basis for the Section 112 rejection seems to be the Examiner's contention that one of ordinary skill in the art would not be "able to learn (i) what is actually expressed as an automaton A; and (ii) whether the abstract program being computed is the result from combining reduced number of states and an altered version of one branch property f, OR this computed program yields reduced number of states and an altered version of a property."

*Point (i)*

Appellant submits that the specification and figures are entirely clear and address point (i). As clearly shown in step 210 of FIG. 2, and discussed in the specification on page 6, lines 19-22, the program abstraction method 200 is initiated with the program, M, to be abstracted, as well as the branching time property, f, to be preserved and an automaton, A, for f.  
20 Clearly, it is the branching time property, f, that is expressed as an automaton, A. Appellant submits that this is also entirely clear in claims 1 and 14.

*Point (ii)*

Both alternative interpretations suggested by the Examiner as part of point (ii) are inaccurate and inconsistent with the actual claim language in claims 1 and 14.

25 The Examiner and the Appeal Board are again referred to FIG. 2, where the program abstraction method 200 is initiated with the program, M, to be abstracted, as well as the

branching time property,  $f$ , to be preserved. The branching time property,  $f$ , is expressed as an automaton,  $A$ .

The exemplary program abstraction method 200 provides two results, both recited in the third step (or limitation) of claims 1 and 14. First, the exemplary program abstraction method 200 computes an abstract program (referred to herein as  $Q$ ) with a reduced number of states. In addition, the exemplary program abstraction method 200 computes an altered property (referred to herein as  $g$ ). Thus,  $g$  is an "altered version of  $f$ ," as recited in claim 1.

The Examiner's main source of confusion seems to be with the second step of claim 1. The examiner's alternative interpretations are:

A. the abstract program being computed is the result from combining reduced number of states and an altered version of one branch property,  $f$ .

This interpretation is incorrect.  $Q$  is created from  $M$  and the \*original\* property,  $f$ .

B. this computed program yields reduced number of states and an altered version of a property.

This interpretation is incorrect in the sense that the computed program,  $Q$ , does not "yield" property  $g$ . Both  $Q$  and  $g$  are computed independently of one another from the given information.

The second step of claim 1 is described in conjunction with step 220 of FIG. 2, as well as page 6, lines 23-28, where it is described that:

Hereafter, the user-specified abstract domain,  $\bar{S}$ , and a set of left-total abstraction relations  $\{\xi_u | u \in Q\}$  are obtained during step 220, where each  $\xi_u \subseteq S \times \bar{S}$ . The abstract domain is a set of values to generalize the possible real values of the program. The abstraction relations relate the real program states to the abstract domain.  $\xi_u$  and  $\xi_v$  are defined to be  $S \times \bar{S}$ , and note that  $(s, q)$  is related to  $(t, q)$  if and only if  $s \xi_q t$  holds.

This passage clearly defines both the abstract domain and abstract relations, as recited in the second step of claim 1, in a completely consistent manner with the claim terminology.

Appellant respectfully requests withdrawal of the Section 112 rejection.

5        Prior Art Rejections: Independent Claims 1 and 14

Independent claims 1 and 14 were rejected under 35 U.S.C. §102(b) as being anticipated by BenAri. Regarding claim 1, the Examiner asserts that BenAri teaches a method for reducing a program, M, that preserves (citing page 5) at least one branching time property, f, comprising the steps of: forming a product of said program, M and said branching time property, f (citing middle of page 208), expressed as an automaton, A (citing bottom of page 208 and top of page 209); obtaining an abstract domain containing a set of abstract values to generalize possible states of said program and abstract relations that relate said program states to said abstract domain (citing sec. 2, page 209, bottom, and Semantics for UB, page 210); and computing an abstract program with a reduced number of states (citing pages 209-224) and an altered version of said branching time property, f (citing pages 212-214), using said product (citing page 210, top, and page 211, top)

BenAri considers whether there exists a program for which a formula is true. BenAri takes a formula of a branching time property and determines if it is true.

20        The present invention, on the other hand, given a program and a property, determines if the property is true of this program.

Among other important limitations found in claims 1 and 14, BenAri does not disclose or suggest *reducing* a program. It is noteworthy that the Examiner has not alleged that this is shown by **any** portion of BenAri.

25        In addition, BenAri does not disclose or suggest *preserving* a branch time property. Rather, BenAri is given a branch time property and identifies a program for which it holds.

BenAri also does not disclose or suggest forming a product of the program, M and

said branching time property, f. It is again noteworthy that the Examiner has not alleged that **any** portion of BenAri describes a product.

BenAri additionally does not disclose an “*abstract domain* containing a set of abstract values to generalize possible states of said program.” The Examiner is merely applying  
5 keyword spotting in this portion of the rejection.

BenAri also does not disclose or suggest “computing an abstract program with a reduced number of states.” The Examiner has not stated how the formulas, axioms, theorems and lemmas read on the elimination of paths.

BenAri also does not disclose or suggest computing “an altered version of said  
10 branching time property, f.” The Examiner has made no attempt to indicate how the branching time property referenced in the first step is altered to form the altered version referenced in the final step.

Thus, other than matching a few key words from the claims of the present invention, BenAri is only peripherally related to the subject matter of the present invention.

15 BenAri does **not** address the issue of ***reducing a program while preserving branching time properties***. In addition, BenAri does not disclose or suggest, among other important limitations, computing an abstract program with a reduced number of states and an altered version of said branching time property, f, using said product, as required by independent claims 1 and 14.

20 Conclusion

The rejections of the cited claims under sections 102 and 103 in view of BenAri are therefore believed to be improper and should be withdrawn. The remaining rejected dependent claims are believed allowable for at least the reasons identified above with respect to the independent claims.

The attention of the Examiner and the Appeal Board to this matter is appreciated.

Respectfully,



Date: February 29, 2008

Kevin M. Mason  
Attorney for Applicant(s)  
Reg. No. 36,597  
Ryan, Mason & Lewis, LLP  
1300 Post Road, Suite 205  
Fairfield, CT 06824  
(203) 255-6560



APPENDIX

1           A method for reducing a program, M, that preserves at least one branching time property, f, comprising the steps of:

5           forming a product of said program, M and said branching time property, f, expressed as an automaton, A;

          obtaining an abstract domain containing a set of abstract values to generalize possible states of said program and abstract relations that relate said program states to said abstract domain; and

10          computing an abstract program with a reduced number of states and an altered version of said branching time property, f, using said product.

2           The method of claim 1, further comprising the step of performing an automated program check

15          3           The method of claim 2, wherein said automated program check is a model checking step.

4           The method of claim 3, wherein said automated program check is performed for  
20          an altered branching time property.

5           The method of claim 1, wherein said computing step further comprises the step of defining a set of states,  $S'$ , in said abstract program as  $S' = \bar{S} \times Q$ , where  $S$  is a set of states in said program, M, and Q is a set of states of the automaton, A.

25          6           The method of claim 5, wherein OR states in said set of states,  $S'$ , are those states where  $\delta(q, true)$  has the form  $q_1 \vee q_2$  or  $\langle a \rangle q_1$ , and all other states are AND states, where q are

individual states and  $\delta$  is a transition relation between states.

7. The method of claim 5, wherein an abstract state  $(t, \hat{q})$  is in a subset of initial states,  $I'$ , of the abstract program if there exists  $s \in I$  for which  $s \xrightarrow{\hat{q}} t$ , where  $s$  is an individual state,  $I$  is a subset of initial states,  $I$ , of the program,  $M$ , and  $\hat{q}$  is one of said abstract relations.

8. The method of claim 5, wherein for an abstract AND state  $(t, q)$ , the transition  $((t, q); (t', q'))$  is in an abstract transition relation,  $R'$ , if there exists a concrete state  $(s, q)$  and a successor  $(s', q')$  that are related to  $(t, q); (t', q')$  respectively.

9. The method of claim 5, wherein for an abstract OR state  $(t, q)$ , the transition  $((t, q); (t', q'))$  is in an abstract transition relation,  $R'$ , only if for every  $(s, q)$  which is related to  $(t, q)$ , there exists a successor  $(s', q')$  which is related to  $(t', q')$ .

10. The method of claim 8, wherein said product  $ATS\ M \times A$  is abstracted by weakening said transition relations at AND states.

11. The method of claim 9, wherein said product  $ATS\ M \times A$  is abstracted by strengthening said transition relations at OR states.

12. The method of claim 8, further comprising the step of obtaining one or more rank functions and employing said one or more rank functions in an abstract transition relation,  $R'$ .

13. The method of claim 8, further comprising the step of obtaining one or more choice predicates and employing said one or more rank functions in an abstract transition

relation,  $R'$ .

14. A system for reducing a program, M, that preserves at least one branching time property, f, comprising:

5 a memory; and  
a processor operatively coupled to said memory, said processor configured to:  
form a product of said program, M and said branching time property, f, expressed  
as an automaton, A;

10 obtain an abstract domain containing a set of abstract values to generalize possible  
states of said program and abstract relations that relate said program states to said abstract  
domain; and

compute an abstract program with a reduced number of states and an altered  
version of said branching time property, f, using said product

15 15. The system of claim 14, wherein said processor is further configured to perform  
an automated program check.

16. The system of claim 15, wherein said automated program check is a model  
checking step.

20

17. The system of claim 16, wherein said automated program check is performed for  
an altered branching time property.

18. The system of claim 14, wherein said processor is further configured to define a  
25 set of states,  $S'$ , in said abstract program as  $S' = \bar{S} \times Q$ , where  $S$  is a set of states in said  
program, M, and Q is a set of states of the automaton, A.

19. The system of claim 18, wherein OR states in said set of states,  $S'$ , are those states where  $\delta(q, true)$  has the form  $q_1 \vee q_2$  or  $\langle a \rangle q_1$ , and all other states are AND states, where  $q$  are individual states and  $\delta$  is a transition relation between states

20. The system of claim 18, wherein an abstract state  $(t, \hat{q})$  is in a subset of initial states,  $I'$ , of the abstract program if there exists  $s \in I$  for which  $s \xi_{\hat{q}}^t$ , where  $s$  is an individual state,  $I$  is a subset of initial states,  $I$ , of the program,  $M$ , and  $\xi_{\hat{q}}$  is one of said abstract relations.

21. The system of claim 18, wherein for an abstract AND state  $(t, q)$ , the transition  $((t, q); (t', q'))$  is in an abstract transition relation,  $R'$ , if there exists a concrete state  $(s, q)$  and a successor  $(s', q')$  that are related to  $(t, q); (t', q')$  respectively.

22. The system of claim 18, wherein for an abstract OR state  $(t, q)$ , the transition  $((t, q); (t', q'))$  is in an abstract transition relation,  $R'$ , only if for every  $(s, q)$  which is related to  $(t, q)$ , there exists a successor  $(s', q')$  which is related to  $(t', q')$ .

23. The system of claim 21, wherein said product  $ATS \ M \times A$  is abstracted by weakening said transition relations at AND states

24. The system of claim 22, wherein said product  $ATS \ M \times A$  is abstracted by strengthening said transition relations at OR states.

25. The system of claim 21, further comprising the step of obtaining one or more rank functions and employing said one or more rank functions in an abstract transition relation,  $R'$ .

26. The system of claim 21, further comprising the step of obtaining one or more choice predicates and employing said one or more rank functions in an abstract transition relation,  $R'$

EVIDENCE APPENDIX

There is no evidence submitted pursuant to § 1.130, 1.131, or 1.132 or entered by the Examiner and relied upon by appellant.

RELATED PROCEEDINGS APPENDIX

There are no known decisions rendered by a court or the Board in any proceeding identified pursuant to paragraph (c)(1)(ii) of 37 CFR 41.37.